# Fighting the poison: DNSSEC to the rescue Stéphane Bortzmeyer

AFNIC bortzme<mark>yer@nic.fr</mark>

# Fighting the poison: DNSSEC to the rescue

Stéphane Bortzmeyer AFNIC bortzmeyer@nic.fr

amic

2 / 23

Data retrieval on the Internet, via a key, the **domain name**. Provides:

afnic- 3/23

Data retrieval on the Internet, via a key, the **domain name**. Provides:

Stability

afnic-

3 / 23

Data retrieval on the Internet, via a key, the **domain name**. Provides:

- Stability
- Memorisability

afnic

3 / 23

Data retrieval on the Internet, via a key, the **domain name**. Provides:

- Stability
- Memorisability
- Security?



3 / 23

Data retrieval on the Internet, via a key, the **domain name**. Provides:

- Stability
- Memorisability
- Security?

Most common data type retrieved: IP addresses

afnic-

3 / 23

Data retrieval on the Internet, via a key, the **domain name**. Provides:

- Stability
- Memorisability
- Security?

Most common data type retrieved: IP addresses

DNS is a vital part of the Internet infrastructure

3 / 23

#### Tree structure

A network database, organized as a tree.



amic



#### • Authoritative servers (masters and slaves) have a pristine copy of the data



#### Name servers

- Authoritative servers (masters and slaves) have a pristine copy of the data
- Resolvers (or recursors or caches or recursive servers) query the authoritative servers

afnic-

#### Name servers

- Authoritative servers (masters and slaves) have a pristine copy of the data
- Resolvers (or recursors or caches or recursive servers) query the authoritative servers
- There is also a stub resolver (often without a cache) in libraries/applications

### Resolution



amic

6 / 23

#### Threats



amic

#### The biggest threat



amic

8 / 23

afnic-9/23

• Communication between authoritative servers and resolvers is typically with UDP  $\longrightarrow$  no protection against IP spoofing

amic

9 / 23

- Communication between authoritative servers and resolvers is typically with UDP  $\longrightarrow$  no protection against IP spoofing
- $\bullet$  The attacker replies before the legitimate server  $\longrightarrow$  done!

afnic-

- Communication between authoritative servers and resolvers is typically with UDP  $\longrightarrow$  no protection against IP spoofing
- $\bullet\,$  The attacker replies before the legitimate server  $\longrightarrow$  done!
- There are some checks by the resolver: query ID (a small cookie), query name...

amic

9 / 23

- Communication between authoritative servers and resolvers is typically with UDP  $\longrightarrow$  no protection against IP spoofing
- $\bullet\,$  The attacker replies before the legitimate server  $\longrightarrow$  done!
- There are some checks by the resolver: query ID (a small cookie), query name...
- Since the data have a Time-To-Live (TTL), if the attacker loses the race, he has to wait

- Communication between authoritative servers and resolvers is typically with UDP  $\longrightarrow$  no protection against IP spoofing
- $\bullet\,$  The attacker replies before the legitimate server  $\longrightarrow$  done!
- There are some checks by the resolver: query ID (a small cookie), query name...
- Since the data have a Time-To-Live (TTL), if the attacker loses the race, he has to wait
- In 2008, Kaminsky discovered a way to retry the attack immediately. This boosted DNSSEC deployment

# Cryptography 101

- DNSSEC uses asymmetric crypto: a key has a private part and a public part. Algorithms: RSA, ECDSA...
- DNSSEC relies on hashing: we sign hashes, not directly the data. Algorithms: SHA

afnic-

10 / 2

**1** Data protection ( $\neq$  channel protection)

afnic 11/2

- **Q** Data protection ( $\neq$  channel protection)
- Check the authenticity of the data, whatever the relays and caches

afnic

11 / 2

- Data protection ( $\neq$  channel protection)
- Oteck the authenticity of the data, whatever the relays and caches
- Sompatible with existing DNS (same resource record format)

amic

- Data protection ( $\neq$  channel protection)
- Oteck the authenticity of the data, whatever the relays and caches
- Sompatible with existing DNS (same resource record format)
- Confidentiality is out of scope

amic

11/2

#### • Each zone has a key (with a public and a private part)

afnic 12/2

- Sech zone has a key (with a public and a private part)
- Resource records are signed with the private part

afnic-

12 / 2

# **DNSSEC** basics

- Each zone has a key (with a public and a private part)
- Resource records are signed with the private part
- O Authoritative name servers serve the signed data

amic

12 / 2

# **DNSSEC** basics

- Sech zone has a key (with a public and a private part)
- ② Resource records are signed with the private part
- O Authoritative name servers serve the signed data
- Validating resolvers check the signature with the public part

amic

# Keys

```
; v Crypto algorithm
; v
absolight.fr. 7069 IN DNSKEY 257 3 8 (
AwEAAateikCxMCJjIPEQ+hKu9xF0RkUtssOkynR7SoUy
...
VtzH7JEEz2Q31qNTWj430m/Bzi8IDCbbfkOlIhk=
) ; key id = 62795
```

```
• 8 \longrightarrow RSA + SHA-256
```

• Key ID (or key tag): a short identifier for the key

amic

13 / 2

#### Signatures

; An ordinary resource record, here of type AAAA (an IP address) absolight.fr. 75018 IN AAAA 2a01:678:2:100::80

```
; The signature
; v Crypto algorithm
; v
absolight.fr. 75018 IN RRSIG AAAA 8 2 86400 20140709092716 (
20140703041612 55713 absolight.fr.
TKwtxqlKiRY5mOcIkJCmrDQRnlxJB5jAja9qScEgQX0j
...
```

- Signed with key 55713 (not the one seen above)
- Valid from 3 july to 9 july

amic

14 / 2

### Chain of trust

How can we be sure we have the right public key?

```
; v Points towards this key
; v
absolight.fr. 161337 IN DS 62795 8 2 (
5C770C1889D8E27DC2606D8A6F5A9B7CF0F943B1F2B7
A66BCBB8F1EEA62582F2 )
```

- DS = Delegation Signer
- A pointer from the parent zone to the public key of the child zone
- Of course, it is signed

• You'll often see two keys, one signing the key set, one signing the data

afnic 16/2

- You'll often see two keys, one signing the key set, one signing the data
- This is not mandatory: co.uk has only one key

afric- 16/2

- You'll often see two keys, one signing the key set, one signing the data
- This is not mandatory: co.uk has only one key
- They are called KSK (Key Signing Key) and ZSK (Zone Signing Key)

amic

- You'll often see two keys, one signing the key set, one signing the data
- This is not mandatory: co.uk has only one key
- They are called KSK (Key Signing Key) and ZSK (Zone Signing Key)
- The idea is to have different characteristics: for instance a short, fast and often changed ZSK and a stable and long KSK

afnic-

16/2

- You'll often see two keys, one signing the key set, one signing the data
- This is not mandatory: co.uk has only one key
- They are called KSK (Key Signing Key) and ZSK (Zone Signing Key)
- The idea is to have different characteristics: for instance a short, fast and often changed ZSK and a stable and long KSK
- In the example above, 62795 was the KSK and 55713 the ZSK

# DNSviz



afnic-

17 / 2

DNSSEC signs records. When there is no record (non-existing domain name, for instance), what do we sign?



18 / 2

DNSSEC signs records. When there is no record (non-existing domain name, for instance), what do we sign?

• We use NSEC or NSEC3 records: they claim "there is nothing here" and are signed for checking

amic

DNSSEC signs records. When there is no record (non-existing domain name, for instance), what do we sign?

- We use NSEC or NSEC3 records: they claim "there is nothing here" and are signed for checking
- NSEC are chained by domain names ("there is nothing between bar.example.org and foo.example.org")

DNSSEC signs records. When there is no record (non-existing domain name, for instance), what do we sign?

- We use NSEC or NSEC3 records: they claim "there is nothing here" and are signed for checking
- NSEC are chained by domain names ("there is nothing between bar.example.org and foo.example.org")
- NSEC3 are chained by hashes of domain names, for more privacy ("there is no domain whose hash is between UI6PC9AJFB1E6GE0GRUL67QNCKIG9BCK and L6M3OP8QM1VR3T47JNM6DBL6S4QM2BL8")

#### How do I do that with free software?

A lot of free programs are available:

- OpenDNSSEC manages the keys life cycle and signs
- For authoritative servers, NSD, Knot, PowerDNS and BIND can serve signed zones
- $\bullet\,$  PowerDNS and BIND can do serving + automatic signatures
- For validating resolvers, Unbound and BIND can check signatures
- To check, Zonecheck, DNScheck, validns...

### Actual deployment

- First TLD signed between 2007 and 2010
- The DNS root was signed in 2010
- Today, all important TLDs are signed
- User domains signed: Internet organizations (ietf.org, afnic.fr...), US federal domains (.gov) or geek domains. No banks or big companies. rmll.info not signed
- Biggest validating resolvers: Google Public DNS and Comcast's DNS service
- $\bullet$  Percentage of protected users: >50 % in Sweden, 25 % in the US, <10 % in France

#### Daily chores

afnic- 21/2



• Monitoring, specially the signatures expiration

afnic- 21/2

# Daily chores

- Monitoring, specially the signatures expiration
- Re-signing: can be done automatically (OpenDNSSEC, for instance)

afnic

21 / 2

# Daily chores

- Monitoring, specially the signatures expiration
- Re-signing: can be done automatically (OpenDNSSEC, for instance)
- Debugging when you manage a validating resolver ("fbi.gov does not work!")

amic

21 / 2

afnic- 22/2

You are now convinced and you want to deploy DNSSEC?

Check the security of your data (remember NY Times vs. SEA)

afnic

22 / 2

- Check the security of your data (remember NY Times vs. SEA)
- Check the quality of your DNS setup (name servers but also middleboxes, for instance broken firewalls limiting data size to 512 bytes)

afnic-

- Check the security of your data (remember NY Times vs. SEA)
- Check the quality of your DNS setup (name servers but also middleboxes, for instance broken firewalls limiting data size to 512 bytes)
- Solution: ONSSEC depends on it

afnic-

- Check the security of your data (remember NY Times vs. SEA)
- Check the quality of your DNS setup (name servers but also middleboxes, for instance broken firewalls limiting data size to 512 bytes)
- Solution: ONSSEC depends on it
- Check the monitoring

afnic-

- Check the security of your data (remember NY Times vs. SEA)
- Check the quality of your DNS setup (name servers but also middleboxes, for instance broken firewalls limiting data size to 512 bytes)
- Solution: ONSSEC depends on it
- Check the monitoring
- (Authoritative service) Think about private key security
- (Authoritative service) Start with a not-too-important zone

amic

You are now convinced and you want to deploy DNSSEC?

- Check the security of your data (remember NY Times vs. SEA)
- Check the quality of your DNS setup (name servers but also middleboxes, for instance broken firewalls limiting data size to 512 bytes)
- Solution: ONSSEC depends on it
- Check the monitoring
- (Authoritative service) Think about private key security
- (Authoritative service) Start with a not-too-important zone
- (Recursive service) Be ready to handle the case of an important zone messing up with DNSSEC

amic

22 / 2



#### Plan in advance: deploying DNSSEC takes time Don't wait the last minute: attackers progress!

afnic

23 / 2

# Merci!

anic

www.afnic.fr contact@afnic.fr

